

# Scalar Quantization with Noisy Partitions and its Application to Flash ADC Design

---

**Da Wang**, Yury Polyanskiy and Gregory Wornell

Signals, Information  
and Algorithms  
Laboratory



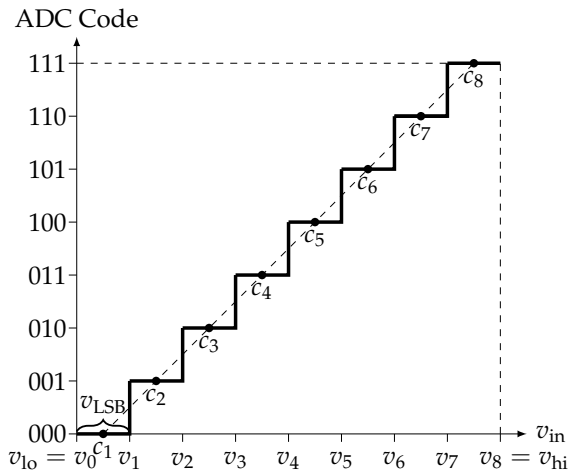
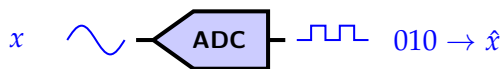
---

*ISIT 2014, Honolulu, HI*

June 30, 2014

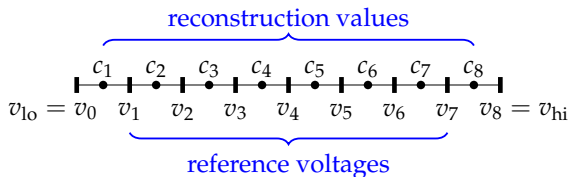
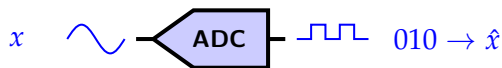
- 1 Background on ADC
  - ▶ Flash ADC architecture
  - ▶ The issue of imprecise comparators
- 2 Scalar Quantization with Noisy Partitions
- 3 High resolution analysis
- 4 ADC design implications

# Analog-to-Digital Converter (ADC)



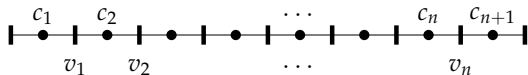
- $2^b$  reconstruction values
- $n = 2^b - 1$  reference voltages

# Analog-to-Digital Converter (ADC)

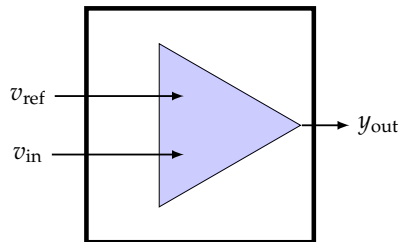


- $2^b$  reconstruction values
- $n = 2^b - 1$  reference voltages

# ADC and its key building block: comparator

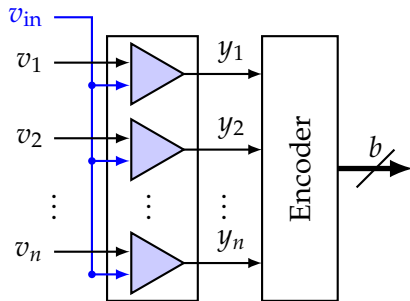


## Comparator



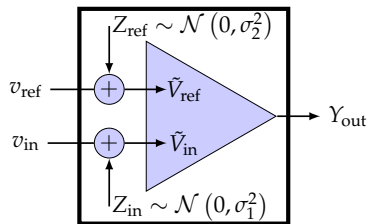
$$y_{\text{out}} = \begin{cases} 1 & v_{\text{in}} > v_{\text{ref}} \\ 0 & v_{\text{in}} \leq v_{\text{ref}} \end{cases}$$

## The Flash ADC architecture



$$n = 2^b - 1$$

# The imprecise comparator due to process variation



$Z_{\text{in}}$  and  $Z_{\text{ref}}$ :

- offsets due to **process variation**
- variation  $\nearrow$  as comparator size  
 $\searrow$
- independent, zero-mean  
Gaussian distributed [Kinget 2005,  
Nuzzo 2008]

**Note:**

- **fixed** after fabrication
- randomness: over a collection of comparators
- aggregate variation:

$$Z = Z_{\text{ref}} - Z_{\text{in}} \sim \mathcal{N}(0, \sigma^2)$$

# A call for mathematical framework

Existing theoretical error analysis (e.g., [Lundin 2005])

- assumes small process variation
- does not attempt to change the design

ADC design with imprecise comparators

- Practice
- ADC with redundancy [Flynn *et al.*, 2003]
  - ADC with redundancy, calibration and reconfiguration [Daly *et al.*, 2008]

# A call for mathematical framework

Existing theoretical error analysis (e.g., [Lundin 2005])

- assumes small process variation
- does not attempt to change the design

ADC design with imprecise comparators

Practice ■ ADC with redundancy [Flynn *et al.*, 2003]

- ADC with redundancy, calibration and reconfiguration [Daly *et al.*, 2008]

Theory ■ Little prior work

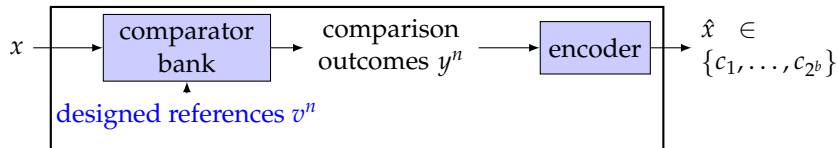
- Related: scalar quantizer with random thresholds for uniform input [Goyal 2011]



# System model: Scalar Quantization with Noisy Partition Points

$b$ -bit ADC

classical

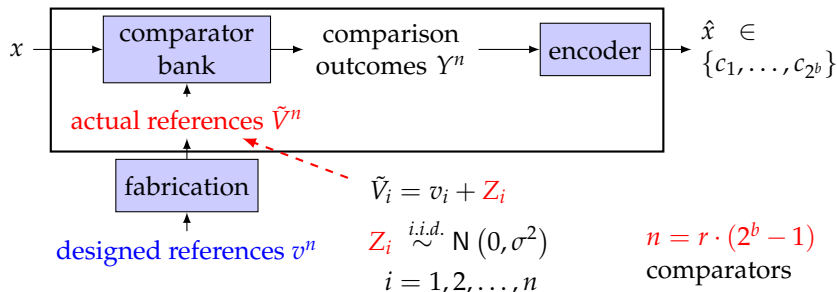


$n = 2^b - 1$   
comparators

# System model: Scalar Quantization with Noisy Partition Points

$b$ -bit ADC

with redundancy

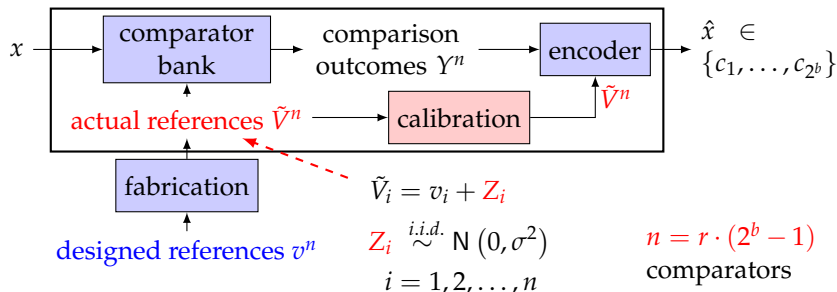


- “references” = “partition points”
- $r$ : redundancy factor

# System model: Scalar Quantization with Noisy Partition Points

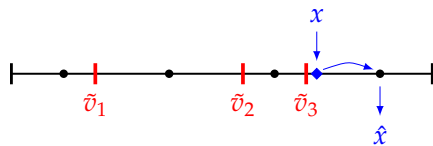
$b$ -bit ADC

with redundancy and calibration



- “references” = “partition points”
- $r$ : redundancy factor

# Performance measures of ADC



error function

$$e(x) = x - \hat{x}$$

mean-square error

$$\text{MSE} = \mathbb{E}_{X, \tilde{V}^n} [e(X)^2]$$

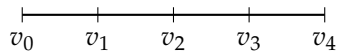
Given input distribution  $f_X$ , how to design **optimal**  $v_1, v_2, \dots, v_n$ ?

$v^n$  —  $\tilde{V}^n$  — analyze MSE

Is **scaling down** the size of comparators **actually beneficial**?

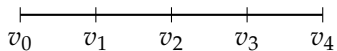
## Challenge: randomness in partition points

e.g., design:

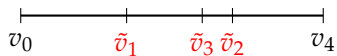


## Challenge: randomness in partition points

e.g., design:

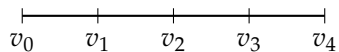


After fabrication:

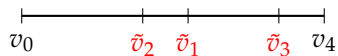
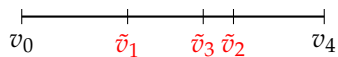


## Challenge: randomness in partition points

e.g., design:



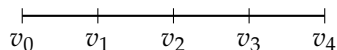
After fabrication:



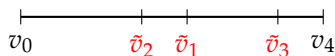
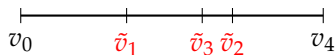
or ...

# Challenge: randomness in partition points

e.g., design:



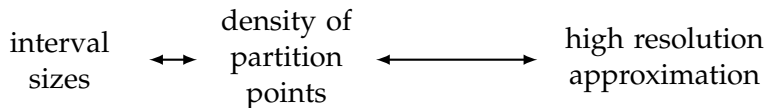
After fabrication:



or ...

## Observations

- Ordering may change  $\rightarrow$  order statistics
- Random interval sizes  $\rightarrow$  ?



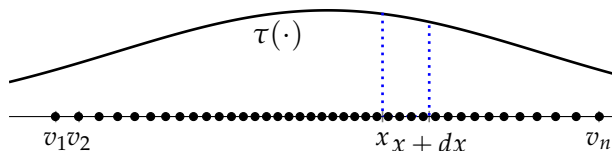


# High resolution approximation

Assume  $n \rightarrow \infty$

- Represent  $v^n$  by **point density functions**  $\tau(x)$

$$\tau(x) dx \approx \frac{\text{number of } v^n \text{ in } [x, x + dx]}{n}$$

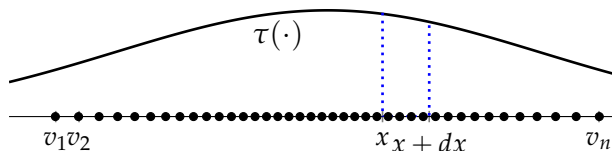


# High resolution approximation

Assume  $n \rightarrow \infty$

- Represent  $v^n$  by **point density functions**  $\tau(x)$

$$\tau(x) dx \approx \frac{\text{number of } v^n \text{ in } [x, x + dx]}{n}$$

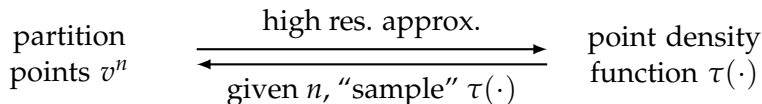


- $\tilde{V}^n$ : **point density functions**  $\lambda(x)$

$$\lambda(x) dx \approx \frac{\mathbb{E} [\text{number of } \tilde{V}^n \text{ in } [x, x + dx]]}{n}$$

- Point density function **simplifies** analysis!

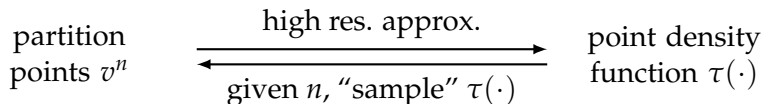
# Point density function guides partition point design



## Examples

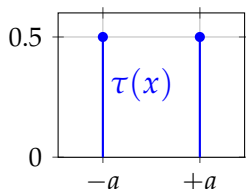
- $\tau \sim \text{Unif}([-1, 1])$
- $v^n$ :  $n$ -point evenly-spaced grid on  $[-1, 1]$

# Point density function guides partition point design

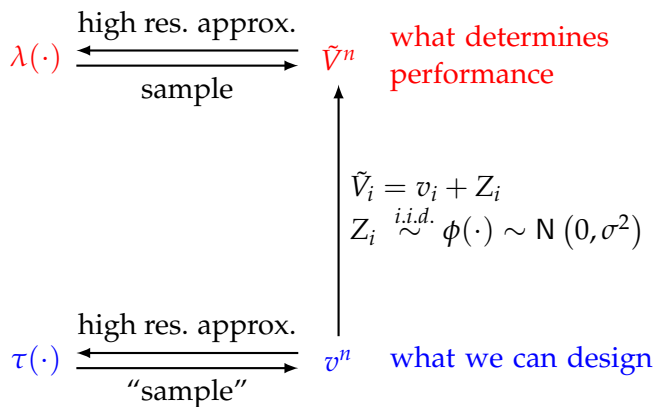


## Examples

- $\tau(x) = 0.5 \cdot \delta(x - a) + 0.5 \cdot \delta(x + a)$
- $v^n$ :
  - ▶  $n/2$  points at  $+a$
  - ▶  $n/2$  points at  $-a$

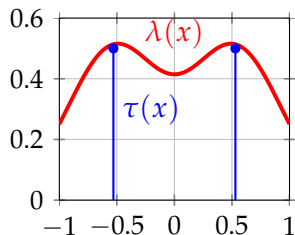
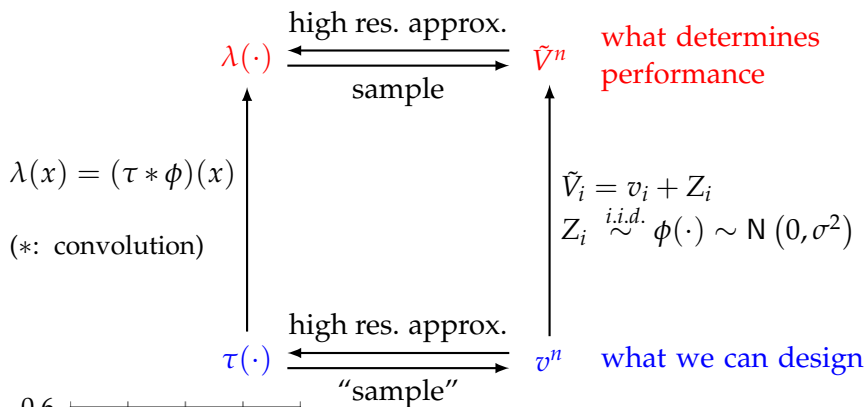


## With process variation, fabricated references matters



- Performance characterization in  $\lambda(\cdot)$
- Want to find the optimal  $\tau(\cdot)$

# With process variation, fabricated references matters



- Performance characterization in  $\lambda(\cdot)$
- Want to find the optimal  $\tau(\cdot)$

## Process variation increases MSE 6-fold

Input  $X \sim f_X(\cdot)$ ,

$$\text{MSE} = \mathbb{E}_X [e(X)^2]$$

classical case [Bennett 1948, Panter & Dite 1951]

$$\text{MSE} \simeq \frac{1}{12n^2} \int \frac{f_X(x)}{\lambda^2(x)} dx$$

$$\lambda = \tau$$

with process variations

$$\text{MSE} \simeq \frac{1}{2n^2} \int \frac{f_X(x)}{\lambda^2(x)} dx$$

$$\lambda = \tau * \phi$$

Why 6 times?

deterministic grid vs. random division of an interval  
(a topic in order statistics)

Optimal  $\tau$

■ a **necessary and sufficient** condition

## Optimal partition point density

Key function:

$$R(\tau) = \int f_X(x)(\tau * \phi)^{-2}(x) dx$$

### Theorem

$\tau$  minimizes  $R(\tau)$  *if and only if*

$$\sup_{x \in \mathcal{A}} \left[ \frac{f_X}{(\tau * \phi)^3} * \phi \right] (x) \leq \left\langle f_X, \frac{1}{(\tau * \phi)^2} \right\rangle.$$

*In particular, if there exists  $\tau^*$  such that*

$$\tau^* * \phi \propto f_X^{1/3},$$

*then  $\tau^*$  minimizes  $R(\tau)$  and*

$$R(\tau^*) = \left( \int f_X^{1/3}(x) dx \right)^3.$$



# MSE-optimal designs can be quite different

## Gaussian input distribution

### Complete characterization of optimal $\tau$

When

$$f_X \sim N(0, \sigma_X^2),$$

then

$$\tau^* \sim \begin{cases} N(0, 3\sigma_X^2 - \sigma^2) & \text{when } 3\sigma_X^2 > \sigma^2 \\ \delta(x) & \text{when } 3\sigma_X^2 \leq \sigma^2 \end{cases}$$

and

$$R(\tau^*) = \begin{cases} 6\sqrt{3}\pi\sigma_X^2 & \text{when } 3\sigma_X^2 > \sigma^2 \\ 2\pi\sigma^3 / \sqrt{\sigma^2 - 2\sigma_X^2} & \text{when } 3\sigma_X^2 \leq \sigma^2 \end{cases}$$

# MSE-optimal designs can be quite different

## Uniform input distribution

$$f_X \sim \text{Unif}([-1, 1])$$

$$\sigma_0 \approx 0.7228$$

$$\sigma < \sigma_0$$

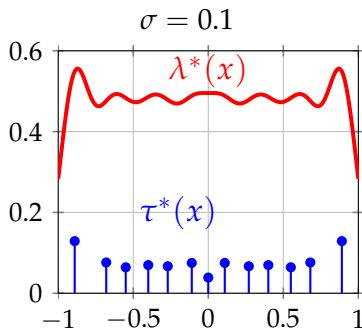
iterative optimization

$\Rightarrow$  **locally optimal**  $\tau^*(x)$

$$\sigma \geq \sigma_0$$

the **necessary and sufficient**  
condition

$\Rightarrow \tau^*(x) = \delta(x)$



# MSE-optimal designs can be quite different

## Uniform input distribution

$$f_X \sim \text{Unif}([-1, 1])$$

$$\sigma_0 \approx 0.7228$$

$$\sigma < \sigma_0$$

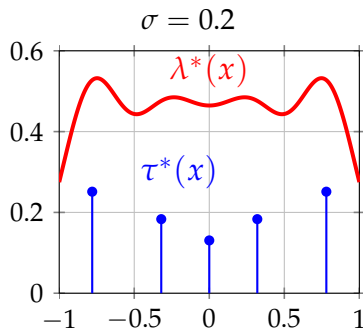
iterative optimization

$\Rightarrow$  locally optimal  $\tau^*(x)$

$$\sigma \geq \sigma_0$$

the necessary and sufficient condition

$\Rightarrow \tau^*(x) = \delta(x)$



# MSE-optimal designs can be quite different

## Uniform input distribution

$$f_X \sim \text{Unif}([-1, 1])$$

$$\sigma_0 \approx 0.7228$$

$$\sigma < \sigma_0$$

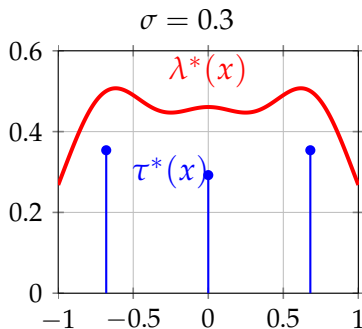
iterative optimization

$\Rightarrow$  **locally optimal**  $\tau^*(x)$

$$\sigma \geq \sigma_0$$

the **necessary and sufficient**  
condition

$\Rightarrow \tau^*(x) = \delta(x)$



# MSE-optimal designs can be quite different

## Uniform input distribution

$$f_X \sim \text{Unif}([-1, 1])$$

$$\sigma_0 \approx 0.7228$$

$$\sigma < \sigma_0$$

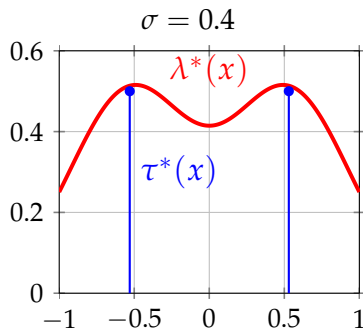
iterative optimization

$\Rightarrow$  locally optimal  $\tau^*(x)$

$$\sigma \geq \sigma_0$$

the necessary and sufficient condition

$\Rightarrow \tau^*(x) = \delta(x)$



# MSE-optimal designs can be quite different

## Uniform input distribution

$$f_X \sim \text{Unif}([-1, 1])$$

$$\sigma_0 \approx 0.7228$$

$$\sigma < \sigma_0$$

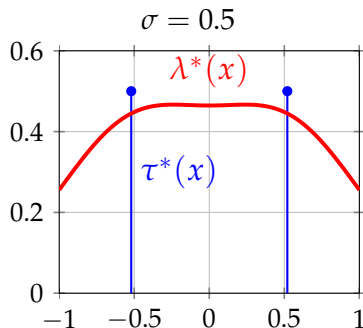
iterative optimization

$$\Rightarrow \text{locally optimal } \tau^*(x)$$

$$\sigma \geq \sigma_0$$

the **necessary and sufficient**  
condition

$$\Rightarrow \tau^*(x) = \delta(x)$$



# MSE-optimal designs can be quite different

## Uniform input distribution

$$f_X \sim \text{Unif}([-1, 1])$$

$$\sigma_0 \approx 0.7228$$

$$\sigma < \sigma_0$$

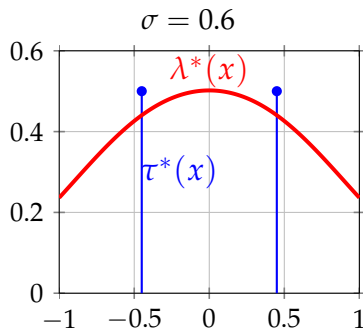
iterative optimization

$\Rightarrow$  **locally optimal**  $\tau^*(x)$

$$\sigma \geq \sigma_0$$

the **necessary and sufficient**  
condition

$\Rightarrow \tau^*(x) = \delta(x)$



# MSE-optimal designs can be quite different

Uniform input distribution

$$f_X \sim \text{Unif}([-1, 1])$$

$$\sigma_0 \approx 0.7228$$

$$\sigma < \sigma_0$$

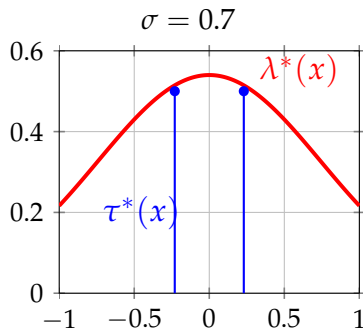
iterative optimization

$\Rightarrow$  locally optimal  $\tau^*(x)$

$$\sigma \geq \sigma_0$$

the necessary and sufficient condition

$\Rightarrow \tau^*(x) = \delta(x)$





# MSE-optimal designs can be quite different

## Uniform input distribution

$$f_X \sim \text{Unif}([-1, 1])$$

$$\sigma_0 \approx 0.7228$$

$$\sigma < \sigma_0$$

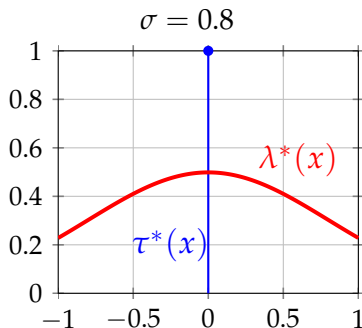
iterative optimization

$\Rightarrow$  locally optimal  $\tau^*(x)$

$$\sigma \geq \sigma_0$$

the necessary and sufficient condition

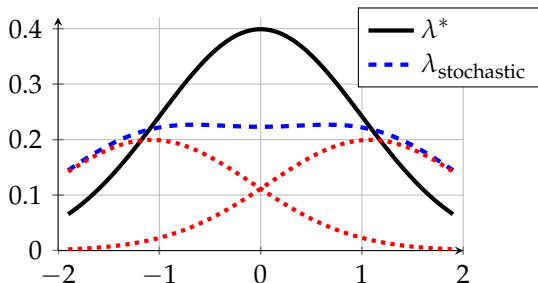
$\Rightarrow \tau^*(x) = \delta(x)$



# ADC design implications

Stochastic ADC [Weaver 2010] is suboptimal

- Uniform input distribution over  $[-1, 1]$ ,  $\sigma = 1.0$



Weaver:

$$\tau(x) = \delta(x - 1.078\sigma) + \delta(x + 1.078\sigma)$$

- Flatter  $\lambda(x)$ , but larger MSE
- Many points out of input range

Optimal:  $\tau(x) = \delta(x)$

- Minimum MSE

$$\text{MSE}_{\text{stochastic}} / \text{MSE}^* \approx 2.15!$$

## Scaling down the size of comparators is beneficial

For circuit fabrication [Kinget 2005, Nuzzo 2008],

$$\text{process variation} \quad \sigma^2 \propto \frac{1}{\text{component area}}$$

Given a fixed silicon area,

$$\text{\# components} \quad n \propto \frac{1}{\text{component area}}$$

Uniform input distribution, when  $\sigma \geq \sigma_0$ ,

$$\text{MSE} \approx 2\pi\sigma^2/n^2 \quad \xrightarrow{\sigma^2 \propto n} \quad \text{MSE} = \Theta(1/n)$$

## Scaling down the size of comparators is beneficial

For circuit fabrication [Kinget 2005, Nuzzo 2008],

$$\text{process variation } \sigma^2 \propto \frac{1}{\text{component area}}$$

Given a fixed silicon area,

$$\text{\# components } n \propto \frac{1}{\text{component area}}$$

Uniform input distribution, when  $\sigma \geq \sigma_0$ ,

$$\text{MSE} \approx 2\pi\sigma^2/n^2 \quad \xrightarrow{\sigma^2 \propto n} \quad \text{MSE} = \Theta(1/n)$$

Building an ADC with more **smaller** but **less precise** comparators improves accuracy!

### Recap

- Scalar quantization with noisy partitions
- High resolution analysis of MSE
- Optimal partition point designs difference from the classical case

### Work in progress

- More error metrics: maximum quantization error, ...
- Partial-calibration or no-calibration
- Take power consumption of ADC into account